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**Method And Apparatus For Protecting Copy Control Information Provided To
A Video Recording Device**

Related Application

5 This application is a continuation-in-part application to U.S. Patent Applications number 09/385,590 and 09/385,592, both entitled Digital Video Content Transmission Ciphering and Deciphering Method and Apparatus, filed on August 29, 1999.

10 BACKGROUND OF THE INVENTION

1. Field of the Invention

 The present invention relates to the field of content protection. More specifically, the present invention addresses the generation of pseudo random
15 numbers for use in a symmetric ciphering/deciphering process based authentication process for authenticating video receiving devices.

2. Background Information

20 ^{sub}_A In general, entertainment, education, art, and so forth (hereinafter collectively referred to as "content") packaged in digital form offer higher audio and video quality than their analog counterparts. However, content producers, especially those in the entertainment industry, are still reluctant in totally embracing the digital form. The primary reason being digital contents are particularly vulnerable to pirating. As
25 unlike the analog form, where some amount of quality degradation generally occurs with each copying, a pirated copy of digital content is virtually as good as the "gold master". As a result, much effort have been spent by the industry in developing and

adopting techniques to provide protection to the distribution and rendering of digital content.

Historically, the communication interface between a video source device (such as a personal computer) and a video sink device (such as a monitor) is an analog interface. Thus, very little focus has been given to providing protection for the transmission between the source and sink devices. With advances in integrated circuit and other related technologies, a new type of digital interface between video source and sink devices is emerging. The availability of this type of new digital interface presents yet another new challenge to protecting digital video content.

While in general, there is a large body of cipher technology known, the operating characteristics such as the volume of the data, its streaming nature, the bit rate and so forth, as well as the location of intelligence, typically in the source device and not the sink device, present a unique set of challenges, requiring a new and novel solution. Parent applications number 09/385,590 and 09/385,592 disclosed various protocol and cipher/deciphering techniques to authenticate a video sink device and protect transmission to the video sink device. Pseudo random numbers are employed as seed or basis numbers for the ciphering/deciphering process.

As technology advances, it is desired to selectively allow certain video sink devices to make authorized copies of the protected video. According, a method and apparatus to protect the provision of copy control information is desired.

BRIEF DESCRIPTION OF DRAWINGS

The present invention will be described by way of exemplary embodiments, but not limitations, illustrated in the accompanying drawings in which like references
5 denote similar elements, and in which:

Figure 1 illustrates an example video source device and an example video sink device incorporated with the teachings of the present invention, in accordance with one embodiment;

Figure 2 illustrates an overview of the method of the present invention for
10 protecting video provided by video source device to video recording device, in accordance with one embodiment;

Figures 3a-3b illustrate the symmetric ciphering/deciphering process in further detail, in accordance with one embodiment;

***Figure 4** illustrates an overview of the method of the present invention for
15 protecting copy control information provided by video source device to video recording device, in accordance with one embodiment;*

Figure 5 illustrates an example combined block/stream cipher unit of **Fig. 1** in further detail, in accordance with one embodiment;

Figure 6 illustrates the block key section of **Fig. 4** in further detail, in
20 accordance with one embodiment; and

Figure 7 illustrates the block data section of **Fig. 4** in further detail, in accordance with one embodiment.

DETAILED DESCRIPTION OF THE INVENTION

In the following description, various aspects of the present invention will be described, and various details will be set forth in order to provide a thorough understanding of the present invention. However, it will be apparent to those skilled in the art that the present invention may be practiced with only some or all aspects of the present invention, and the present invention may be practiced without the specific details. In other instances, well known features are omitted or simplified in order not to obscure the present invention.

Various operations will be described as multiple discrete steps performed in turn in a manner that is most helpful in understanding the present invention. However, the order of description should not be construed as to imply that these operations are necessarily performed in the order they are presented, or even order dependent. Lastly, repeated usage of the phrase "in one embodiment" does not necessarily refer to the same embodiment, although it may.

*Referring now to **Figure 1**, wherein a block diagram illustrating an example video source device incorporated with the teachings of the present invention, in accordance with one embodiment, is shown. As illustrated, video source device **102**, incorporated with the teachings of the present invention, and video sink or recording device **104** are coupled to each other via digital video link **106**. Video source device **102** includes authentication unit **108** and video hardware interface **110**, sharing cipher unit **112**. While not detailed, video sink/recording device **104** is also similarly constituted with its own authentication unit, hardware interface and cipher unit.*

*Video source device **102**, using authentication unit **108**, authenticates video sink/recording device **104**. Authentication unit **108** employs an authentication*

process that is based on a symmetric ciphering/deciphering process, requiring a seed or basis value to be provided by authentication unit **108**. The seed or basis value is a pseudo random number. Upon authenticating video sink/recording device **104**, video source device **102** generates and provides video content to video sink/recording device **104** through video hardware interface **110**, which ciphers video content before transmission to protect the video content from unauthorized copying. Video hardware interface **110** ciphers the video content using ciphering bits generated by cipher unit **112**. In accordance with the present invention, video source device **102** also provides copy control information to video sink/recording device **104** to specifically authorize and control copying or recording of the video content provided. Further, video source device **102** also protects these copy control information to prevent their tampering. As will be described in more detail below, video source device **102** advantageously protects these copy control information by coupling or tying them to the symmetric ciphering and deciphering process employed to protect the video content. As a result, the copy control information can not be tampered with, as tampering with the copy control information will cause video sink/recording device **104** to be unable to decipher the ciphered video.

The exact nature of these copy control information is application dependent and of no particular relevance to the practice of the present invention. Similarly, the communication interface employed (not shown) as well as the communication protocol employed by video source device **102** to convey the copy control information to video sink/recording device is also of no particular relevance to the practice of the present invention. Any communication link and protocol known in the art may be employed.

Except for the teachings of the present invention incorporated, to be described more fully below, video source device **102** is intended to represent a

broad range of digital devices known in the art, including but not limited to computers of all sizes (from palm size device to desktop device, and beyond), set-up boxes, or DVD players. Examples of video recording devices include but are not limited to computing devices with storage medium, "digital VCR" and the like. As to

5 *digital video link 106, it may be implemented in any one of a number of mechanical and electrical forms, as long as they are consistent with the operating requirement (i.e. speed, bit rate and so forth), and a mechanism (which may be in hardware or through protocol) is provided to allow control information to be exchanged between video source and sink/recording devices 102 and 104.*

10 *Before proceeding to describe the present invention in further detail, it should be noted that video sink/recording device 104 may also be disposed "behind" a video signal repeater device, repeating signals for the "remotely" disposed video sink/recording device 104, as opposed to being directly coupled to video source device 102 as illustrated.*

15 **Figure 2** illustrates an overview of the symmetric ciphering/deciphering process based method for providing video content from a source device to a sink/recording device, in accordance with one embodiment. In this embodiment, source and sink/recording devices **102** and **104** are assumed to have each been

20 provided with an array of private keys and a complementary identifier by a certification authority. As illustrated, upon power on or reset, source device **102** first provides a basis value to the symmetric ciphering/deciphering process to sink/device device **104** (block **202**). For the illustrated embodiment, the basis value is a random number (A_n). A_n may be generated in any one of a number of

25 techniques known in the art. Additionally, source device **102** also provides a selected one of its device keys ($A_{k_{sv}}$) to sink/recording device **104** (block **202**). In

response, sink/recording device **104** replies with a selected one of its device keys (Bk_{sv}) (block **203**). Upon exchanging the above information, source and sink/recording devices **102** and **104** independently generate their respective copies of an authentication key (K_m) using Ak_{sv} and Bk_{sv} (block **204** and **205**). For the illustrated embodiment, source device **102** generates its copy of K_m by summing private keys of its provided array indexed by Bk_{sv} , while sink/recorder device **104** generates its copy of K_m by summing private keys of its provided array indexed by Ak_{sv} . At this time, if both source and sink devices **102** and **104** are authorized devices, they both possess and share a common secret authentication key K_m .

In one embodiment, each of source and sink/recording devices **102** and **104** is pre-provided with an array of 40 56-bit private keys by the certification authority. A_n is a 64-bit random number, and K_m is 56-bit long. For more information on the above described authentication process, see co-pending U.S. Patent Application, serial number 09/275,722, filed on March 24, 1999, entitled Method and Apparatus for the Generation of Cryptographic Keys, having common inventorship as well as assignee with the present application.

Having authenticated sink/recording device **104**, source device **102** ciphers video content into a ciphered form before transmitting the video content to sink device **104**. Source device **102** ciphers the video content employing a symmetric ciphering/deciphering process, and using the random number (A_n) as well as the independently generated authentication key (K_m) (block **206**). Upon receipt of the video content in ciphered form, sink/recording device **104** decipheres the ciphered video content employing the same symmetric ciphering/deciphering processing, and using the provided A_n as well as its independently generated copy of K_m (block **207**).

In accordance with the present invention, as an integral part of ciphering video content, source device **102** derives a set of verification reference values in a

predetermined manner (block **208**). Likewise, as an integral part of symmetrically deciphering video content sink/recording device **104** also derives a set of verification values in a predetermined manner, and transmits these derived verification values to source device **102** (block **209**). Upon receiving each of these verification values,

5 source device **102** compares the received verification value to the corresponding one of the verification reference values to determine and confirm that indeed the ciphered video content is being properly deciphered by sink/recording device **104** (block **210**).

For the illustrated embodiment, both source and sink/recording devices **102**

10 and **104** generate the verification reference and verification values continuously, but the verification values are provided from sink/recording device **104** to source device **102** periodically at predetermined intervals.

In one embodiment, the verification reference and verification values are all 64-bits in length, and sink/recording device **104** provides source device **102** with

15 verification values at initialization and every 64th frames thereafter.

Figures 3a-3b illustrate the symmetric ciphering/deciphering process in further detail, in accordance with one embodiment. In this embodiment, the video content is assumed to be a multi-frame video content with each frame having

20 multiple lines of video content. In between two lines of a frame is an interval to allow a sink device to horizontally "retrace" itself, commonly known as the horizontal retrace interval or horizontal blanking interval (HBI). Likewise, in between two frames is an interval to allow a sink device to vertically "retrace" itself, commonly known as the vertical retrace interval or vertical blanking interval (VBI).

25 Source device **102** first generates a session key (K_s) for the transmission session (block **302**). For the illustrated embodiment, K_s is generated by block

ciphering the above mentioned random number A_n using the authentication key K_m as the block cipher key and applying C1 clocks. The duration of a transmission session is application dependent. Typically, it corresponds to a natural demarcation of the video content, e.g. the transmission of a single movie may constitute a transmission session, or the transmission of an episode of a sitcom may constitute a transmission session instead.

Upon generating the session key K_s , source device **102** generates an initial version of a second random number (M_0) (block **304**). For the illustrated embodiment, source device **102** first generates a pseudo random bit sequence (at p-bit per clock) using a stream cipher with the above described random number A_n and the session key K_s (in two roles, as another input random number and as the stream cipher key), applying C2 clocks. Source device **102** derives M_0 from the pseudo random bit sequence, as the bit sequence is generated.

Next, source device **102** generates a frame key (K_i) for the next frame (block **306**). For the illustrated embodiment, K_i is generated by block ciphering an immediately preceding version of the second random number M_{i-1} using the session key K_s as the block cipher key, and applying C3 clocks. That is, for the first frame, frame-1, frame key K_1 is generated by block ciphering the above described initial version of the second random number M_0 , using K_s , and applying C3 clocks.

Additionally, this operation is subsequently repeated at each vertical blanking interval for the then next frame, frame-2, frame-3, and so forth.

Upon generating the frame key K_i , source device **102** generates the current version of the second random number (M_i) (block **302**). For the illustrated embodiment, source device **102** first generates a pseudo random bit sequence (at p-bit per clock) using a stream cipher with the previous version of the second random number M_{i-1} and the frame key K_i (in two roles, as another input random

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for protecting the video content. In one embodiment, video source device **102** incorporates the *n*-bit copy control information as an integral part of an initialization value to be employed to initialize a basis register of a round function of the block cipher section of cipher unit **112**. More specifically, in one embodiment, video
5 source device **102** incorporates the *n*-bit copy control information as the most significant bits (MSB) of the initialization value for the data section round function of the block cipher.

Upon receipt of the copy control information from video source device **102**, video sink/recording device **104** would do the same thing, that is incorporating the *n*-
10 bit copy control information into the corresponding initialization value for a corresponding register of its cipher unit in like manner.

Thereafter, **404**, video source and sink/recording device **102** and **104** would correspondingly initialize their cipher units, including the special initialization value (incorporated with the copy control information) among the initialization values to be
15 employed. Accordingly, the copy control information affects the pseudo random sequence subsequently generated by cipher unit **112**, and employed to cipher video content as earlier described. In like manner, unless tampered, the corresponding effects will manifest themselves in the pseudo random sequence subsequently generated by the cipher unit of video sink/recording device **104** to decipher the
20 ciphered video content.

In the event the copy control information are tampered with, the correct pseudo random sequence required to decipher the ciphered video content will not be produced, and as a result, the video content will not be available for copy.

Thus, it can be seen from the foregoing, the copy control information are
25 protected from tampering.

Figure 5 illustrates an example combined block/stream cipher suitable for use to practice the present invention, in accordance with one embodiment. As illustrated, combined block/stream cipher **112** includes block key section **502**, data section **504**, stream key section **506**, and mapping section **508**, coupled to one another as shown. Block key section **502** and data section **504** are employed in both the block mode as well as the stream mode of operation, whereas stream key section **506** and mapping section **508** are employed only in the stream mode of operation.

Briefly, in block mode, block key section **502** is provided with a block cipher key, such as the earlier described authentication key K_m or the session key K_s ; whereas data section **504** is provided with the plain text, such as the earlier described random number A_n or the derived random number M_{i-1} . "Rekeying enable" signal is set to a "disabled" state, operatively de-coupling block key section **502** from stream key section **506**. During each clock cycle, the block cipher key as well as the plain text are transformed. The block cipher key is independently transformed, whereas transformation of the plain text is dependent on the transformation being performed on the block cipher key. After a desired number of clock cycles, the provided plain text is transformed into ciphered text. For the earlier described video content protection method, when block key section **502** is provided with K_m and data section **504** is provided with the A_n , ciphered A_n is read out and used as the session key s . When block key section **502** is provided with K_s and data section **504** is provided with the M_{i-1} , ciphered M_{i-1} is read out and used as the frame key K_i .

To decipher the ciphered plain text, block key section **502** and data section **504** are used in like manner as described above to generate the intermediate "keys", which are stored away (in storage locations not shown). The stored

5 embodiment each, referencing **Figs. 6-7.**

0 “Rekeying enable” signal is set to an “enabled” state, operatively coupling block key section **502** to stream key section **506**. Periodically, at predetermined intervals, such as the earlier described horizontal blanking intervals, stream key section **506** is used to generate one or more data bits to dynamically modify the then current state of the random number stored in block data section **502**. During each clock cycle, in
5 between the predetermined intervals, both random numbers stored in block key section **502** and data section **504** are transformed. The random number provided to block key section **502** is independently transformed, whereas transformation of the random number provided to data section **504** is dependent on the transformation being performed in block key section **502**. Mapping block **506** retrieves a subset
20 each, of the newly transformed states of the two random numbers, and reduces them to generate one bit of the pseudo random bit sequence. Thus, in a desired number of clock cycles, a pseudo random bit sequence of a desired length is generated.

For the illustrated embodiment, by virtue of the employment of the “rekeying enable” signal, stream key section **506** may be left operating even during the block

mode, as its outputs are effectively discarded by the “rekeying enable” signal (set in a “disabled” state).

Figure 6 illustrates the block key section of **Fig. 5** in further detail, in accordance with one embodiment. As illustrated, block key section **502** includes registers **602a-602c**, substitution boxes **604**, and linear transformation unit **606**. In block mode, registers **602a-602c** are collectively initialized to a block cipher key, e.g. authentication key K_m or session key K_s . In stream mode, registers **602a-602c** are collectively initialized to a random number, e.g. session key K_s or frame key K_f . Each round, substitution boxes **604** and linear transformation unit **606** modify the content of registers **602a-602c**. More specifically, substitution boxes **604** receive the content of register **602a**, modify it, and then store the substituted content into register **602c**. Similarly, linear transformation unit **606** receives the content of registers **602b** and **602c**, linearly transforms them, and then correspondingly stores the linearly transformed content into registers **602a** and **602b**.

Substitution boxes **604** and linear transformation unit **606** may be implemented in a variety of ways in accordance with well known cryptographic principles. One specific implementation is given in more detail below after the description of **Fig. 7**.

Figure 7 illustrates the block data section of **Fig. 5** in further detail, in accordance with one embodiment. For the illustrated embodiment, data section **504** is similarly constituted as block key section **502**, except linear transformation unit **706** also takes into consideration the content of register **602b**, when transforming the contents of registers **702b-702c**. In block mode, registers **702a-702c** are collectively initialized with the target plain text, e.g. earlier described random number

A_n or derived random number M_{i-1} (incorporated with the n-bit copy control information). In stream mode, registers **702a-702c** are collectively initialized with a random number. Each round, substitution boxes **704** and linear transformation unit **706** modify the content of registers **702a-702c** as described earlier for block key section **502** except for the differences noted above.

Again, substitution boxes **604** and linear transformation unit **606** may be implemented in a variety of ways in accordance with well known cryptographic principles.

In one implementation for the above described embodiment, each register **602a, 602b, 602c, 702a, 702b, 702c** is 28-bit wide. [Whenever registers **602a-602c** or **702a-702c** collectively initialized with a key value or random number less than 84 bits, the less than 84-bit number is initialized to the lower order bit positions with the higher order bit positions zero filled.] Additionally, each set of substitution boxes **604** or **704** are constituted with seven 4 input by 4 output substitution boxes. Each linear transformation unit **606** or **706** produces 56 output values by combining outputs from eight diffusion networks (each producing seven outputs). More specifically, the operation of substitution boxes **604/704** and linear transformation unit **606/706** are specified by the four tables to follow. For substitution boxes **604/704**, the l th input to box J is bit $l*7+J$ of register **602a/702a**, and output l of box J goes to bit $l*7+j$ of register **602c/702c**. [Bit 0 is the least significant bit.] For each diffusion network (linear transformation unit **606** as well as **706**), the inputs are generally labeled $I0-I6$ and the outputs are labeled $O0-O6$. The extra inputs for each diffusion network of the linear transformation unit **706** is labeled $K0-K6$.

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
SK0	8	14	5	9	3	0	12	6	1	11	15	2	4	7	10	13
SK1	1	6	4	15	8	3	11	5	10	0	9	12	7	13	14	2
SK2	13	11	8	6	7	4	2	15	1	12	14	0	10	3	9	5
SK3	0	14	11	7	12	3	2	13	15	4	8	1	9	10	5	6
SK4	12	7	15	8	11	14	1	4	6	10	3	5	0	9	13	2
SK5	1	12	7	2	8	3	4	14	11	5	0	15	13	6	10	9
SK6	10	7	6	1	0	14	3	13	12	9	11	2	15	5	4	8
SB0	12	9	3	0	11	5	13	6	2	4	14	7	8	15	1	10
SB1	3	8	14	1	5	2	11	13	10	4	9	7	6	15	12	0
SB2	7	4	1	10	11	13	14	3	12	15	6	0	2	8	9	5
SB3	6	3	1	4	10	12	15	2	5	14	11	8	9	7	0	13
SB4	3	6	15	12	4	1	9	2	5	8	10	7	11	13	0	14
SB5	11	14	6	8	5	2	12	7	1	4	15	3	10	13	9	0
SB6	1	11	7	4	2	5	12	9	13	6	8	15	14	0	3	10

Table I – Substitution performed by each of the seven constituting substitution boxes of substitution boxes **604/704**.

0:0:0:0:0:0:0:0:0:0

	Diffusion Network Logic Function
O₀	$K_0 \oplus I_1 \oplus I_2 \oplus I_3 \oplus I_4 \oplus I_5 \oplus I_6$
O₁	$K_1 \oplus I_0 \oplus I_2 \oplus I_3 \oplus I_4 \oplus I_5 \oplus I_6$
O₂	$K_2 \oplus I_0 \oplus I_1 \oplus I_3 \oplus I_4 \oplus I_5 \oplus I_6$
O₃	$K_3 \oplus I_0 \oplus I_1 \oplus I_2 \oplus I_4 \oplus I_5 \oplus I_6$
O₄	$K_4 \oplus I_0 \oplus I_1 \oplus I_2 \oplus I_3 \oplus I_5 \oplus I_6$
O₅	$K_5 \oplus I_0 \oplus I_1 \oplus I_2 \oplus I_3 \oplus I_4 \oplus I_6$
O₆	$K_6 \oplus I_0 \oplus I_1 \oplus I_2 \oplus I_3 \oplus I_4 \oplus I_5 \oplus I_6$

	K1	K2	K3	K4	K5	K6	K7	K8
I₀	Kz0	Ky0	Ky4	Ky8	Ky12	Ky16	Ky20	Ky24
I₁	Kz1	Ky1	Ky5	Ky9	Ky13	Ky17	Ky21	Ky25
I₂	Kz2	Ky2	Ky6	Ky10	Ky14	Ky18	Ky22	Ky26
I₃	Kz3	Ky3	Ky7	Ky11	Ky15	Ky19	Ky23	Ky27
I₄	Kz4	Kz7	Kz10	Kz13	Kz16	Kz19	Kz22	Kz25
I₅	Kz5	Kz8	Kz11	Kz14	Kz17	Kz20	Kz23	Kz26
I₆	Kz6	Kz9	Kz12	Kz15	Kz18	Kz21	Kz24	Kz27
O₀	Kx0	Ky0	Ky1	Ky2	Ky3	Kx7	Kx8	Kx9
O₁	Kx1	Ky4	Ky5	Ky6	Ky7	Kx10	Kx11	Kx12
O₂	Kx2	Ky8	Ky9	Ky10	Ky11	Kx13	Kx14	Kx15
O₃	Kx3	Ky12	Ky13	Ky14	Ky15	Kx16	Kx17	Kx18
O₄	Kx4	Ky16	Ky17	Ky18	Ky19	Kx19	Kx20	Kx21
O₅	Kx5	Ky20	Ky21	Ky22	Ky23	Kx22	Kx23	Kx24
O₆	Kx6	Ky24	Ky25	Ky26	Ky27	Kx25	Kx26	Kx27

Tables II & III – Diffusion networks for linear transformation unit 606/706
(continued in Table IV).

	B1	B2	B3	B4	B5	B6	B7	B8
I₀	Bz0	By0	By4	By8	By12	By16	By20	By24
I₁	Bz1	By1	By5	By9	By13	By17	By21	By25
I₂	Bz2	By2	By6	By10	By14	By18	By22	By26
I₃	Bz3	By3	By7	By11	By15	By19	By23	By27
I₄	Bz4	Bz7	Bz10	Bz13	Bz16	Bz19	Bz22	Bz25
I₅	Bz5	Bz8	Bz11	Bz14	Bz17	Bz20	Bz23	Bz26
I₆	Bz6	Bz9	Bz12	Bz15	Bz18	Bz21	Bz24	Bz27
K₀	Ky0	–	–	–	–	Ky7	Ky14	Ky21
K₁	Ky1	–	–	–	–	Ky8	Ky15	Ky22
K₂	Ky2	–	–	–	–	Ky9	Ky16	Ky23
K₃	Ky3	–	–	–	–	Ky10	Ky17	Ky24
K₄	Ky4	–	–	–	–	Ky11	Ky18	Ky25
K₅	Ky5	–	–	–	–	Ky12	Ky19	Ky26
K₆	Ky6	–	–	–	–	Ky13	Ky20	Ky27
O₀	Bx0	By0	By1	By2	By3	Bx7	Bx8	Bx9
O₁	Bx1	By4	By5	By6	By7	Bx10	Bx11	Bx12
O₂	Bx2	By8	By9	By10	By11	Bx13	Bx14	Bx15
O₃	Bx3	By12	By13	By14	By15	Bx16	Bx17	Bx18
O₄	Bx4	By16	By17	By18	By19	Bx19	Bx20	Bx21
O₅	Bx5	By20	By21	By22	By23	Bx22	Bx23	Bx24
O₆	Bx6	By24	By25	By26	By27	Bx25	Bx26	Bx27

Table IV – Diffusion networks for linear transformation unit **606/706** (continued from Tables II & III).

- 5 Referring now back to **Fig. 5**, recall that a ciphered text may be deciphered by generating the intermediate “keys” and applying them backward. Alternatively, for an embodiment where either the inverse of substitution boxes **604/704** and linear transformation units **606/706** are included or they may be dynamically reconfigured to operate in an inverse manner, the ciphered text may be deciphered as follows.
- 10 First, the cipher key used to cipher the plain text is loaded into block key section **502**, and block key section **502** is advanced by R-1 rounds, i.e. one round short of

the number of rounds (R) applied to cipher the plain text. After the initial R-1 rounds, the ciphered text is loaded into data section **504**, and both sections, block key section **502** and data section **504**, are operated "backward", i.e. with substitution boxes **604/704** and linear transformation units **606/706** applying the inverse
5 substitutions and linear transformations respectively.

Other sections, such as stream key section **506** and mapping section **508**, of the example combined block/stream cipher illustrated in **Fig. 5**, are of no particular significance to the practice of the present invention. Accordingly, they will not be "re-described" in detail here. These sections are described in detail in the parent
10 applications.

Accordingly, a novel method and apparatus for protecting copy control information provided to a video recording device has been described.

15 Epilogue

From the foregoing description, those skilled in the art will recognize that many other variations of the present invention are possible. Thus, the present invention is not limited by the details described, instead, the present invention can be practiced with modifications and alterations within the spirit and scope of the
20 appended claims.
